

A journey into the world of p -adic heights

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based on joint projects with

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Koç University

Math Seminar

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- time to time, I'll cheat a bit...

1 Preliminaries

- p -adic numbers
- Elliptic curves and abelian varieties
- Curves and their Jacobians

2 Classical Heights

- Néron–Tate height pairing
- Birch and Swinnerton–Dyer (BSD) conjecture

3 p -adic Heights & Our Results

- Motivation: p -adic BSD and quadratic Chabauty
- Coleman–Gross height pairing
- Mazur–Tate height pairing
- Schneider height pairing

4 Future Work

- A dream: quadratic Chabauty at bad primes
- p -adic BSD

5 Acknowledgements & References

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Example. For $n \in \mathbb{Z}$,

$$v_p(p^n) = n, \quad |p^n|_p = p^{-n}.$$

Then $|p^n|_p$ is small when n is large; $p^n \rightarrow 0$ (p -adically!).

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There is a general philosophy in Number Theory that “all completions are created equal” and should have the same rights. - M. Stoll

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Definition. An **elliptic curve** is a curve of the form

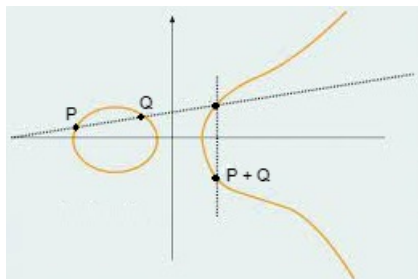
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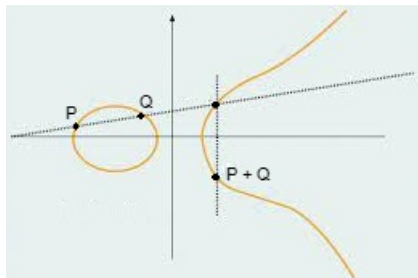


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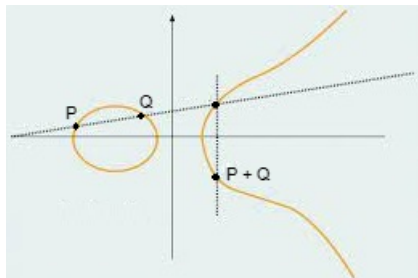
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It is possible to write endlessly on elliptic curves. (This is not a threat.) - S. Lang

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Theorem (Mordell–Weil). For an abelian variety A/\mathbb{Q} , the group $A(\mathbb{Q})$ of rational points of A is a **finitely-generated abelian group**; that is,

$$A(\mathbb{Q}) \simeq \mathbb{Z}^r \oplus A(\mathbb{Q})_{\text{tors}}$$

for some $r \in \mathbb{Z}_{\geq 0}$, called the **algebraic rank** of A/\mathbb{Q} .

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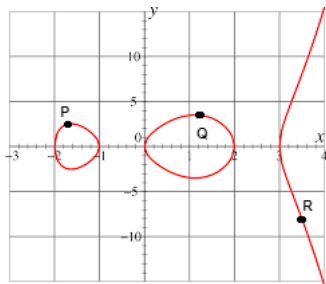
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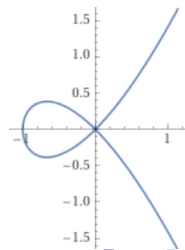
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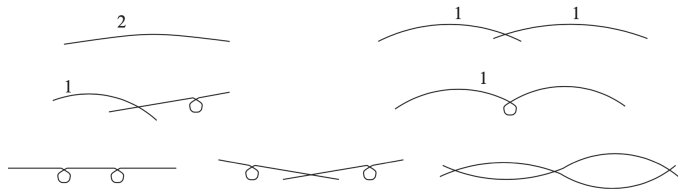
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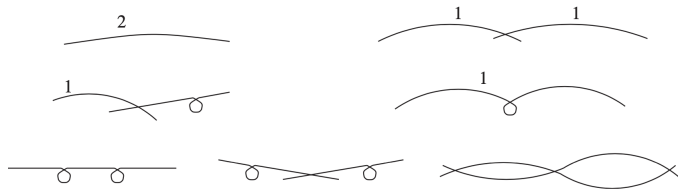
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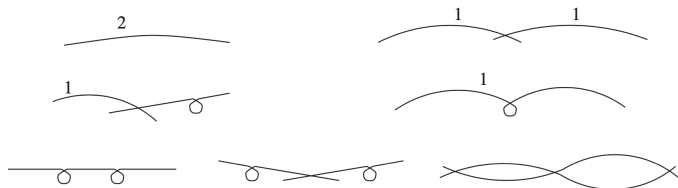
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Algebraic curves were created by God and algebraic surfaces by the Devil.

- M. Noether

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$$\begin{aligned} \iota: C &\hookrightarrow J \\ P &\mapsto P - O \end{aligned}$$

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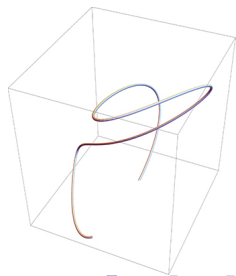
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The variety J is called the **Jacobian variety** of C . For a fixed point $O \in C$, we have an embedding

$$\begin{aligned} \iota: C &\hookrightarrow J \\ P &\mapsto P - O \end{aligned}$$



§2.1. Size of points

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- $h(r_1)$ and $h(r_2)$ are very **far**.

§2.1. Néron–Tate height pairing

The naive height is

$$h_{\text{naive}}: \mathbb{P}^n(\mathbb{Q}) \rightarrow \mathbb{R}, \quad x \mapsto \log \left(\max \{ |x_0|, |x_1|, \dots, |x_n| \} \right)$$

where $x = (x_0 : x_1 : \dots : x_n)$, $x_i \in \mathbb{Z}$ and $\gcd(x_0, x_1, \dots, x_n) = 1$.

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The Néron–Tate height is

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§2.2. BSD conjecture for abelian varieties

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Here, $\text{Reg}(J/\mathbb{Q})$ is the canonical regulator, i.e.,

$$\text{Reg}(J/\mathbb{Q}) = \left| \det \begin{pmatrix} \langle P_1, P_1 \rangle^{\text{NT}} & \langle P_1, P_2 \rangle^{\text{NT}} & \cdots & \langle P_1, P_r \rangle^{\text{NT}} \\ \vdots & \vdots & \ddots & \vdots \\ \langle P_r, P_1 \rangle^{\text{NT}} & \langle P_r, P_2 \rangle^{\text{NT}} & \cdots & \langle P_r, P_r \rangle^{\text{NT}} \end{pmatrix} \right| \in \mathbb{R}$$

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This remarkable conjecture relates the behavior of a function L , at a point where it is not at present known to be defined, to the order of a group III , which is not known to be finite. - J. Tate

§3.1. p -adic heights and motivation

Fix a prime number p . A p -adic height pairing is a function

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- play a crucial role in carrying out the quadratic Chabauty method to determine rational points on curves of genus ≥ 2 .

The p -adic height pairing constructed by Schneider is particularly important because

- the corresponding p -adic regulator fits into p -adic versions of Birch and Swinnerton-Dyer conjecture.

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The local components away from p are described using “arithmetic intersection theory”, and

$$\langle D_1, D_2 \rangle_p^{\text{CG}} := \int_{D_2}^{\text{Vol}} \omega_{D_1}$$

where

- ω_{D_1} is a “canonical” differential form attached to D_1 , and
- \int^{Vol} is the Vologodsky integration.

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§3.2. p -adic integration theories

Example: Consider the hyperelliptic curve X/\mathbb{Q}_5 given by

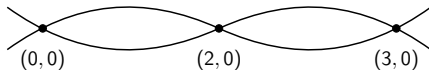
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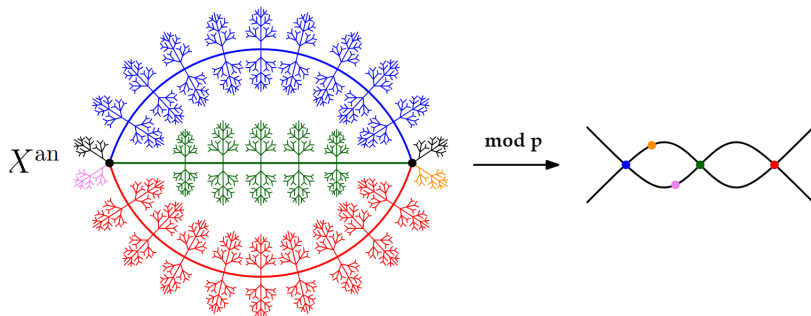
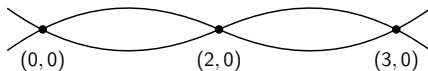


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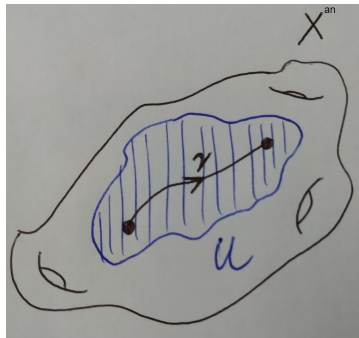
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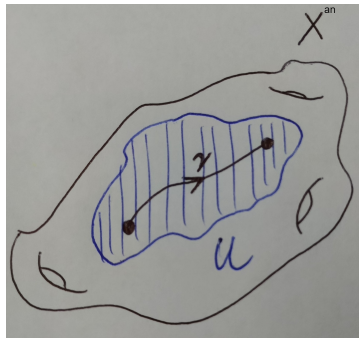
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$$\begin{aligned} \text{BC} \int_\gamma \omega &: \text{BC-integral on } X^{\text{an}} \\ &= \\ \text{BC} \int_\gamma \omega|_U &: \text{BC-integral on } U \end{aligned}$$

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Theorem (Katz–K, K.)

We have

$${}^{\text{Vol}}\int_P \omega = {}^{\text{BC}}\int_{\gamma} \omega - \sum_i \left(c_i \cdot {}^{\text{BC}}\int_{\gamma_i} \omega \right)$$

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Algorithm (Katz–K, K.)

Compute Vologodsky integrals on hyperelliptic curves using this formula and the fact that the Berkovich–Coleman integral is local.

§3.2. Computing Coleman–Gross height pairing

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Determining ω_{D_1} from D_1 is *tricky*...

§3.3. Mazur–Tate height pairing

The pairing

$$\langle \cdot, \cdot \rangle^{\text{MT}} : \text{Div}^0(C) \times \text{Div}^0(C) \rightarrow \mathbb{Q}_p$$

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The real-valued Néron–Tate height pairing $\langle \cdot, \cdot \rangle^{\text{NT}}$ can be decomposed into a sum of “Néron” functions. Here is a p -adic analogue:

§3.3. Mazur–Tate height pairing

Theorem (Bianchi–K.–Müller)

For each q , there exists a p -adic “Néron function” λ_q such that the local pairing $\langle \cdot, \cdot \rangle_q^{\text{MT}}$ can be expressed in terms of λ_q .

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Here is a **direct local** comparison of Coleman–Gross and Mazur–Tate heights in a special case:

Theorem (Bianchi–K.–Müller)

Let C be a hyperelliptic curve of genus 2. For each prime q , we have $\langle \cdot, \cdot \rangle_q^{\text{CG}} = \langle \cdot, \cdot \rangle_q^{\text{MT}}$.

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Theorem (Bianchi–K.–Müller)

For each q , there exists a p -adic “Néron function” λ_q such that the local pairing $\langle \cdot, \cdot \rangle_q^{\text{MT}}$ can be expressed in terms of λ_q .

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Algorithm (Bianchi–K.–Müller)

Compute the Mazur–Tate p -adic height pairing on Jacobians of hyperelliptic curves of genus 2 using the “**theta expression**” of λ_q .

§3.4. Schneider height pairing

The pairing

$$\langle \cdot, \cdot \rangle^{\text{Sch}} : \text{Div}^0(C) \times \text{Div}^0(C) \rightarrow \mathbb{Q}_p$$

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As usual, the local pairing $\langle \cdot, \cdot \rangle_p^{\text{Sch}}$ is more **tricky** to define... But, a nice formula was given by Werner in the case where C/\mathbb{Q}_p is a *Mumford curve*.

§3.4. Mumford curves

Definition. C/\mathbb{Q}_p is called a **Mumford curve** if $C \bmod \mathbb{F}_p$ is a union of curves of genus 0.

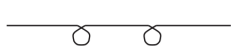
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Assume from now on C/\mathbb{Q}_p is a Mumford curve. Then it admits a **p -adic uniformization**: there exists a p -adic “Schottky” group Γ together with a “ p -adic analytic” isomorphism

$$C \simeq \Omega/\Gamma$$

where $\Omega = \mathbb{P}^1(\mathbb{C}_p) \setminus \mathbb{P}^1(\mathbb{Q}_p)$, the **p -adic upper half plane**.

§3.4. Werner's formula for $\langle \cdot, \cdot \rangle_p^{\text{Sch}}$

Fix two parameters $a, b \in \Omega$, and define the **theta function** on Ω :

$$\Theta(a, b; z) := \prod_{\gamma \in \Gamma} \frac{z - \gamma(a)}{z - \gamma(b)}, \quad z \in \Omega.$$

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Theorem (Werner). Choose preimages x', y', z', w' in Ω . We then have

$$\langle D, E \rangle_p^{\text{Sch}} = \log_p \left(\frac{\Theta(x', y'; z')}{\Theta(x', y'; w')} \right) - \text{another function in terms of } \Theta.$$

§3.4. Computing Schneider height pairing

Algorithm (K.–Masdeu–Müller–van der Put (in progress))

*Compute the Schneider p -adic height pairing on Jacobians of hyperelliptic **Mumford** curves.*

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Theorem (K.–Masdeu–Müller–van der Put)

$$H(z) := \Theta(a, \gamma(a); z) \cdot \prod_{i=0}^g \Theta(a_i, b; z) \cdot \Theta(b_i, s_0(b); z), \quad z \in \Omega$$

is such a function.

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There are no bad primes, really. - N. Dogra

§4.2. p -adic BSD for abelian varieties

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On the other hand, the MTT conjecture in the case of split multiplicative reduction, the “exceptional” case, is of **special interest**. One might expect that a **generalization** of this conjecture to higher dimensional abelian varieties in the case of “split purely toric” reduction can be formulated.



with Masdeu and Müller, Benasque (2022)



with Bianchi and Müller, Groningen (2023)



The journey has ended :-) Teşekkürler!

- *Diophantine Geometry: An Introduction* - Hindry–Silverman
- *Fundamentals of Diophantine Geometry* - Lang

- *p-adic heights on curves* - Coleman–Gross
- *Canonical height pairings via biextensions* - Mazur–Tate
- *p-adic height pairings I* - Schneider
- *Local Heights on Mumford Curves* - Werner
- *Computational tools for quadratic Chabauty* - Balakrishnan–Müller

- *Algorithms for Schneider heights on Mumford curves (in progress)* - K.–Masdeu–Müller–van der Put
- *Algorithms for Coleman–Gross Heights on Hyperelliptic Curves (in preparation)* - Bianchi–K.–Müller
- *Coleman–Gross heights and p-adic Néron functions on Jacobians of genus 2 curves* - Bianchi–K.–Müller
- *Explicit Vologodsky integration for hyperelliptic curves* - K.
- *p-adic Integration on bad reduction hyperelliptic curves* - Katz–K.